

# *(Strict) coherence on Łukasiewicz events: geometry, probability and logic*

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ANCONA

Coherence, strict coherence and MV-algebras

The probability of (strict) coherence

The geometry of coherence: from maps to points

The logic of coherence: from sets to formulas

# A glimpse on (strict) coherence

In 1931 Bruno de Finetti provided a general justification for the probabilistic representation of rational degrees of belief defining the *probability* of an unknown event  $a$  as the *fair price* that a rational Gambler is willing to pay to participate in a betting game, against the Bookmaker, the payoffs of which are 1 in case  $a$  occurs, and 0 otherwise.

Based on this very simple idea, de Finetti showed that all theorems of Kolmogorov's probability theory may be derived as consequences of his *coherence* criterion on assignments (books) on logically connected events.



Consider a finite class of (unknown), classical (yes/no), events  $\Phi = \{e_1, \dots, e_k\}$  and two players: the *Bookmaker*  $\mathcal{B}$  and the Gambler  $\mathcal{G}$ .

- ▶  $\mathcal{B}$  publishes her betting quotients in a *book*, i.e., a map  $\beta : \Phi \rightarrow [0, 1]$ .
- ▶  $\mathcal{G}$  decides stakes  $\sigma_i$  (positive or negative!) and pays  $\sum_{i=1}^k \sigma_i \cdot \beta(e_i)$ .
- ▶ Once a truth-valuation  $h$  is realized,  $\mathcal{B}$  pays back to  $\mathcal{G}$  the amount  $\sum_{i=1}^k \sigma_i \cdot h(e_i)$ .
- ▶ The *balance* for the bookmaker in  $h$  is  $\sum_{i=1}^k \sigma_i(\beta(e_i) - h(e_i))$ .



Given  $\Phi$ , a book  $\beta$  is said to be:

- *Coherent* if for every choice of stakes  $\sigma_1, \dots, \sigma_k$ , there exists at least a possible world  $h$  in which Bookmaker's balance is not negative, i.e.,

$$\sum_{i=1}^k \sigma_i(\beta(e_i) - h(e_i)) \geq 0$$

- *Strictly-coherent* is for every choice of stakes  $\sigma_1, \dots, \sigma_k$ , if there exists a possible world  $h$  in which

$$\sum_{i=1}^k \sigma_i(\beta(e_i) - h(e_i)) < 0$$

then, there must exist another possible world  $h'$  in which

$$\sum_{i=1}^k \sigma_i(\beta(e_i) - h'(e_i)) > 0.$$

MV-algebras are structures in the signature  $\langle \oplus, \neg, 0 \rangle$  of type  $(2, 1, 0)$  which forms the variety generated by the standard algebra:

$$[0, 1]_{MV} = \langle [0, 1], \oplus, \neg, 0 \rangle$$

where

$$a \oplus b = \min\{1, a + b\}, \neg a = 1 - a, 0 = 0.$$

The variety  $\mathbf{MV}$  of MV-algebras is the *equivalent algebraic semantics* of Łukasiewicz infinite-valued logic,  $\mathcal{L}_\infty$ . Thus, in particular, formulas of  $\mathcal{L}_\infty$  are naturally translated into algebraic terms in the language of MV-algebras so that *tautologies* naturally corresponds to *theorem*.

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Further operations are definable as follows

$$a \odot b = \neg(\neg a \oplus \neg b) = \max\{0, a + b - 1\};$$

$$a \rightarrow b = \neg a \oplus b = \min\{1, 1 - a + b\};$$

$$1 = \neg 0;$$

$$a \wedge b = a \odot (a \rightarrow b) = \min\{a, b\};$$

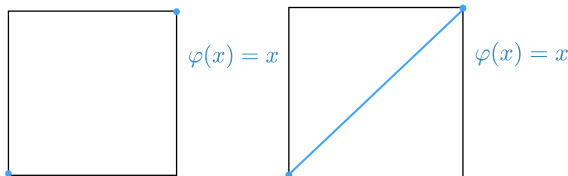
$$a \vee b = \neg(\neg a \wedge \neg b) = \max\{a, b\}.$$



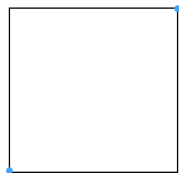
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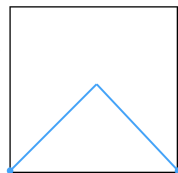
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- (3). The free MV-algebra over  $n$  free generators  $\mathcal{F}_n$ . This is the MV-subalgebra of  $[0, 1]^{[0, 1]^n}$  of those functions which are (i) continuous; (ii) piecewise linear; (iii) each piece has integer coefficient. These are usually called *McNaughton functions*.



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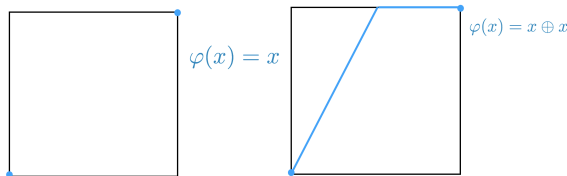


$$\varphi(x) = x$$



$$\varphi(x) = x \wedge \neg x$$

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Fix any finite subset  $\Phi = \{f_1, \dots, f_k\}$  of  $\mathcal{F}_n$ , the free MV-algebra over  $n$ -free generators and let  $\beta : \Phi \rightarrow [0, 1]$  be a book. Then:

1.  $\beta$  is *coherent* if for all  $\sigma_1, \dots, \sigma_k \in \mathbb{R}$  there exists an MV-homomorphism  $h : \mathcal{F}_n \rightarrow [0, 1]_{MV}$  (i.e., a *possible world*) such that

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2.  $\beta$  is *strictly-coherent* if for all  $\sigma_1, \dots, \sigma_k \in \mathbb{R}$ , if there exists an MV-homomorphism  $h : \mathcal{F}_n \rightarrow [0, 1]_{MV}$  (i.e., a *possible world*) such that

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then, there must exist another homomorphism  $h'$  such that

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# The probability of (strict) coherence

## Theorem

Fix a finite subset of classical formulas (i.e., classical events)  $\Phi = \{e_1, \dots, e_k\}$  and a book  $\beta : \Phi \rightarrow [0, 1]$ . Then

1.  $\beta$  is coherent iff it extends to a finitely additive probability measure on the boolean algebra spanned by events in  $\Phi$ . <sup>(a)</sup>
2.  $\beta$  is strictly-coherent iff it extends to a regular and finitely additive probability measure on the boolean algebra spanned by events in  $\Phi$ . <sup>(b and c)</sup>

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<sup>a</sup>B. de Finetti, Theory of Probability, vol. 1, John Wiley and Sons, New York, 1974.

<sup>b</sup>J.G. Kemeny, Fair bets and inductive probabilities, JSL 20(3): 263–273, 1955.

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As for MV-algebras...

### Definition (Mundici, 1995)

By a *state* of an MV-algebra  $\mathbf{A}$  we mean a function  $s : A \rightarrow [0, 1]$  such that  $s(1) = 1$  and,  $s(a \oplus b) = s(a) + s(b)$ , whenever  $a \odot b = 0$ .

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### Proposition (Kroupa, 2005 and Panti, 2009)

A map  $s : \mathcal{F}_n \rightarrow [0, 1]$  is a state iff  $s$  is of the form

$$s(f) = \int_{[0,1]^n} f \, d\mu$$

for  $\mu$  being a Borel regular probability measure on  $[0, 1]^n$ .

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### Theorem (Mundici, 2006)

A book  $\beta$  on a finite subset  $\Phi$  of a free MV-algebra  $\mathcal{F}_n$  is coherent iff there exists a state  $s$  of  $\mathcal{F}_n$  which extends  $\beta$ .

## Definition (Mundici, 1995)

A *faithful state* of an MV-algebra  $\mathbf{A}$  is state  $s : A \rightarrow [0, 1]$  such that  $s(a) = 0$ , implies  $a = 0$ .

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A book  $\beta$  on a finite subset  $\Phi$  of a free MV-algebra  $\mathcal{F}_n$  is strictly-coherent iff there exists a faithful state  $s$  of  $\mathcal{F}_n$  which extends  $\beta$ .

# The geometry of coherence

For every finite subset  $\Phi = \{f_1, \dots, f_k\}$  of a free MV-algebra  $\mathcal{F}_n$ , we shall denote by

$$\mathcal{D}_\Phi = \{\beta : \Phi \rightarrow [0, 1] \mid \beta \text{ is coherent}\}$$

and

$$\mathcal{S}_\Phi = \{\beta : \Phi \rightarrow [0, 1] \mid \beta \text{ is strictly coherent}\}$$

Thus, a book  $\beta : \Phi \rightarrow [0, 1]$ , displayed as  $\beta = \langle \beta(f_1), \dots, \beta(f_k) \rangle$  is a point of  $[0, 1]^k$ . Therefore,  $\mathcal{D}_\Phi$  and  $\mathcal{S}_\Phi$  are subsets of  $[0, 1]^k$  as well.

For each subset  $\Phi = \{f_1, \dots, f_k\}$  of  $\mathcal{F}_n$ , both  $\mathcal{D}_\Phi$  and  $\mathcal{S}_\Phi$  are convex subsets of  $[0, 1]^k$ . (see <sup>a</sup> and <sup>b</sup>)

<sup>a</sup>D. Mundici, Bookmaking on infinite-valued events. Int. J. Approx. Reasoning, 43(3): 223–249, 2006.

<sup>b</sup>T. Flaminio, Three Characterizations of Strict Coherence on Infinite-valued Events. Submitted, 2018

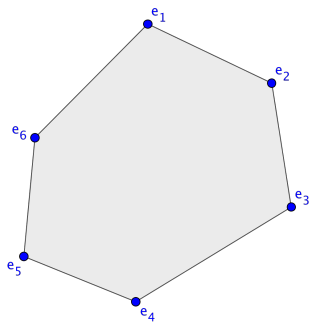
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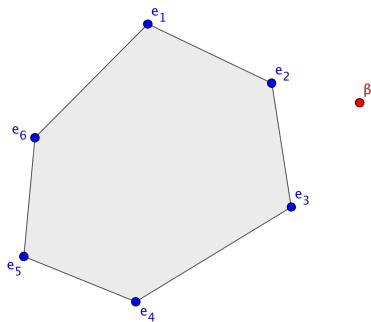
Let  $\Phi$  be a finite subset of  $\mathcal{F}_n$ . Then  $\mathcal{D}_\Phi$  is a polytope and

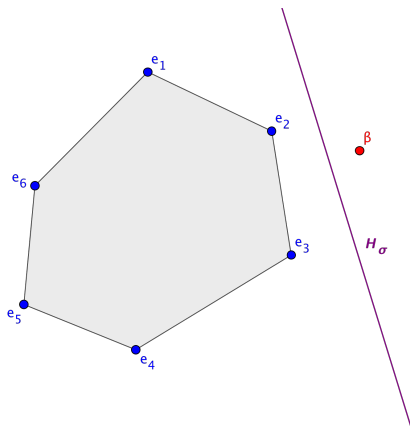
$$\mathcal{S}_\Phi = \text{relint } \mathcal{D}_\Phi.^a$$

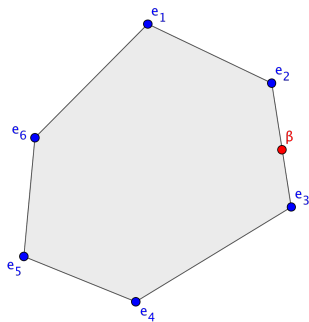
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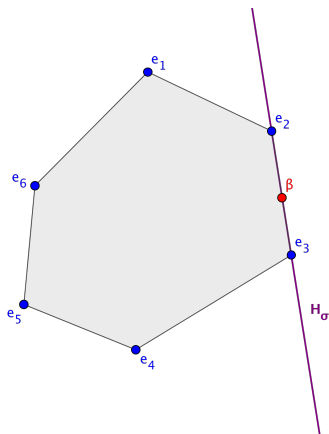


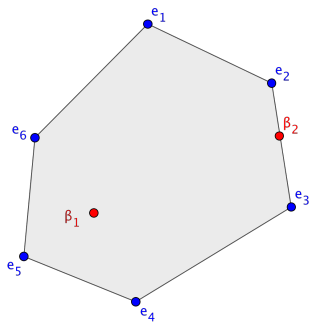


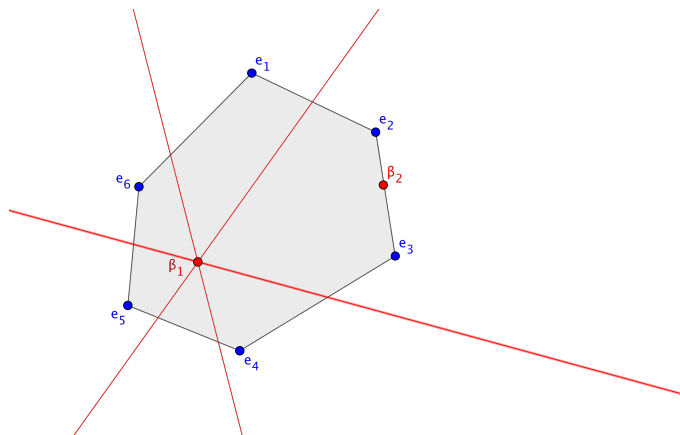


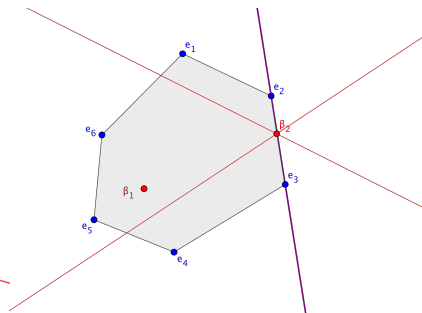
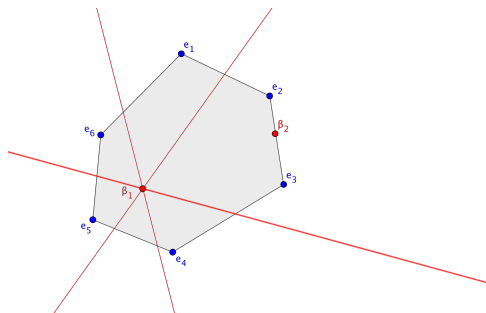












# The logic of coherence (from sets to formulas)

For each rational polyhedron  $P$  of  $[0, 1]^k$ , there exists a formula  $\Pi_P$  of Łukasiewicz logic with  $k$  variables such that

$$P = \mathcal{Mod}(\Pi_P).$$

Furthermore, if  $P_1$  and  $P_2$  are rational polyhedra,

$$P_1 \subseteq P_2 \text{ iff } \mathcal{Mod}(\Pi_{P_1}) \subseteq \mathcal{Mod}(\Pi_{P_2}) \text{ iff } \Pi_{P_1} \models \Pi_{P_2}.$$

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The previous result, due to Mundici, has been further extended by Marra and Spada who provided a categorical duality between the category of *finitely presented MV-algebras* (with homomorphisms) and the category of *rational polyhedra* (with  $\mathbb{Z}$ -maps).



However, it is not difficult to show that, for every rational polyhedra  $P$ , the formula  $\Pi_P$  can be effectively determined.

There exists an effective procedure  $\Pi$  which computes, for each finite  $\Phi \subseteq \mathcal{F}_n$  with  $|\Phi| = k$  and for each  $\beta \in [0, 1]^k$ , Łukasiewicz formulas

$$\Pi_\Phi, \Pi_{(\text{rb } \Phi)} \text{ and } \Pi_\beta$$

The onesets of which are the rational polyhedra  $\mathcal{C}_\Phi$ ,  $\text{rb } \mathcal{C}_\Phi$  and  $\{\beta\}$ .

### Theorem (Flaminio, 2018)

Let  $\Phi = f_1, \dots, f_k$  be a finite set of  $\mathcal{F}_n$  and let  $\beta$  be a rational-valued book on  $\Phi$ . Then the following hold:

1.  $\beta$  is coherent iff  $\vdash \Pi_\beta \rightarrow \Pi_\Phi$ .
2.  $\beta$  is strictly coherent iff  $\vdash \Pi_\beta \rightarrow \Pi_\Phi$  and  $\not\vdash \Pi_\beta \rightarrow \Pi_{(\text{rb } \Phi)}$ .

Thank you.